

# Notes on Beck's *Distributive Laws*

L. Ze Wong

University of Washington, Seattle

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# WARNING!

The notation in this set of notes differs from Beck's paper in the following key ways:

- ▶ Beck writes composites in the opposite direction:  $GF$  means applying  $G$  first, then  $F$ . We will use  $GF$  to mean  $F$  then  $G$ .
- ▶ 'Triple' = 'monad', 'cotriple' = 'comonad'
- ▶ 'Tripleable' = 'monadic', i.e. equivalent to the adjunction involving the category of algebras over monad.

## Motivation 1: Multiplication over Addition

Let  $S$  be the free monoid monad,  $T$  the free abelian group monad.

'Multiplication distributes over addition' means we have a map:

$$STX \rightarrow TSX$$

$$\text{e.g. } (a + b)(c + d) \mapsto ac + ad + bc + bd$$

where  $X = \{a, b, c, \dots\}$ , say.

Further,  $TS$  is the free ring monad.

## Motivation 2: Tensoring monoids

Let  $A, B$  be monoids in a *braided* monoidal category  $(\mathcal{V}, \otimes, 1)$ .

Then  $A \otimes B$  is also a monoid, with multiplication

$$A \otimes B \otimes A \otimes B \xrightarrow{A \otimes tw \otimes B} A \otimes A \otimes B \otimes B \xrightarrow{m_A \otimes m_B} A \otimes B$$

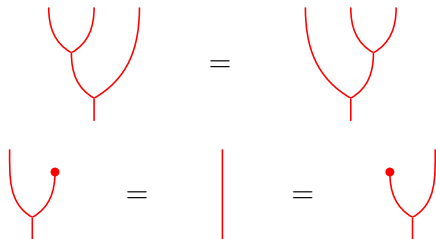
where  $tw : B \otimes A \rightarrow A \otimes B$  is given by the braiding.

# Monads in a 2-category

Fix a 2-category  $\mathbf{K}$ . A *monad* in  $\mathbf{K}$  consists of:

- ▶ 0-cell  $\mathbf{X}$
- ▶ 1-cell  $S : \mathbf{X} \rightarrow \mathbf{X}$
- ▶ 2-cells  $\eta^S : 1_{\mathbf{X}} \Rightarrow S$  and  $\mu^S : SS \Rightarrow S$

such that



i.e. a monad is a monoid in the monoidal category  $(\text{End}(\mathbf{X}), \circ, 1_{\mathbf{X}})$ , for some 0-cell  $\mathbf{X}$ .

# Distributive Law

A distributive law of  $S$  over  $T$  is a 2-cell  $\ell : ST \Rightarrow TS$



such that:



# Characterization of Distributive Laws

# Characterization

Theorem (Beck 1969, Street 1972, Cheng 2011)

*The following are equivalent:*

1. Distributive laws  $\ell : ST \Rightarrow TS$ ,
2. Multiplications  $m : TSTS \Rightarrow TS$  s.t.  $(TS, \eta^T \eta^S, m)$  is monad satisfying the *middle unitary law*, and

$$S \xrightarrow{\eta^T \eta^S} TS \xleftarrow{T \eta^S} T$$

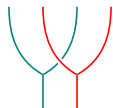
are monad morphisms.

3. *Liftings* of the monad  $T$  to a monad  $\tilde{T}$  over  $\mathbf{X}^S$ ,
4. *Extensions* of the monad  $S$  to a monad  $\tilde{S}$  over  $\mathbf{X}_T$ ,
5. Certain elements of  $\mathbf{Mnd}(\mathbf{Mnd}(\mathbf{K}))$ .

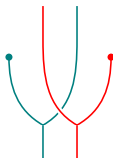


# The composite monad

Given  $\ell : ST \Rightarrow TS$ , define  $m : TSTS \Rightarrow TS$  to be

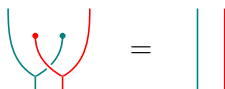


To get back  $\ell$ , do:



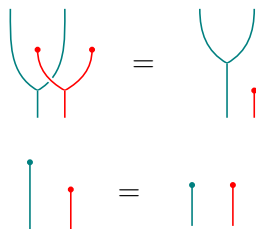
# The composite monad

The *middle unitary law* holds:



The diagram shows an equality between two configurations of lines. On the left, there are two blue lines that curve upwards and meet at a point. From this junction, two red lines curve upwards and meet at a point. On the right, there are two vertical lines, one blue and one red, standing side-by-side.

and  $T\eta^S : T \Rightarrow TS$  is a monad morphism:



The diagram shows two equalities. The first equality shows a configuration of two blue lines and two red lines on the left, where the blue lines meet at a point and the red lines meet at a point, with the blue lines crossing over the red lines. This is equal to a configuration on the right with two blue lines meeting at a point and a single red line to the right. The second equality shows a configuration on the left with a single blue line and a single red line, which is equal to a configuration on the right with a blue line and a red line standing side-by-side.

Similarly,  $\eta^T S : S \Rightarrow TS$  is a monad morphism.

# Liftings and Extensions

A *lift* of  $T$  to the EM object  $\mathbf{X}^S$  is a monad  $\tilde{T}$ :

$$\begin{array}{ccc} & & \mathbf{X}^S \\ & \nearrow \tilde{T} & \downarrow U^S \\ \mathbf{X}^S & \xrightarrow{TU^S} & \mathbf{X} \end{array} \quad + \text{ compatibility equations}$$

An *extension* of  $S$  to the Kleisli object  $\mathbf{X}_T$  is a monad  $\tilde{S}$ :

$$\begin{array}{ccc} \mathbf{X} & \xrightarrow{F_T S} & \mathbf{X}_T \\ F_T \downarrow & \nearrow \tilde{S} & \\ \mathbf{X}_T & & \end{array} \quad + \text{ compatibility equations}$$

Kleisli objects in  $\mathbf{K}$  are EM objects in  $\mathbf{K}^{op}$ , so proofs for liftings hold for extensions too, by duality.

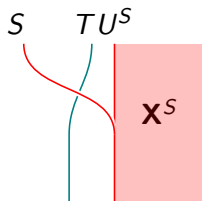
# Liftings and Extensions

Universal property<sup>1</sup> of  $\mathbf{X}^S$ :

$$\{ \text{Functors } \tilde{G} : \mathbf{Y} \rightarrow \mathbf{X}^S \} \cong \left\{ \begin{array}{l} \text{Functors } G : \mathbf{Y} \rightarrow \mathbf{X} \\ \text{with } S\text{-action } \sigma : SG \Rightarrow G \end{array} \right\}$$

$$\begin{array}{ccc} & & \mathbf{X}^S \\ & \nearrow \tilde{G} & \downarrow U^S \\ \mathbf{Y} & \xrightarrow{G} & \mathbf{X} \end{array}$$

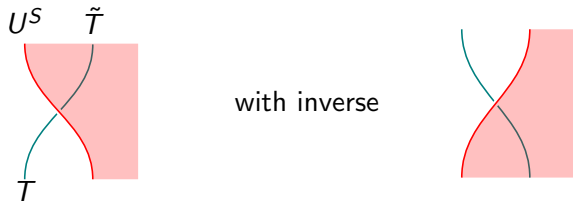
Let  $\mathbf{Y} = \mathbf{X}^S$ ,  $G = TU^S$ ,  $\tilde{T} = \tilde{G}$ . Need  $S$ -action  $STU^S \Rightarrow TU^S$ .  
Given by distributive law and canonical action of  $S$  on  $U^S$ :



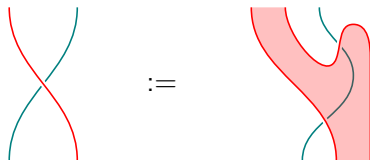
<sup>1</sup>In fact, this is an equivalence of *categories*

# Liftings and Extensions

Conversely, a lifting  $\tilde{T}$  means we have invertible 2-cells:



Lets us define a distributive law:



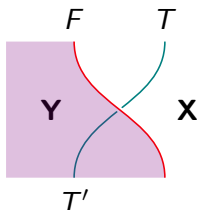
This works for lifts over any adjunction that gives  $S!$

# Monads in $\mathbf{Mnd}(\mathbf{K})$

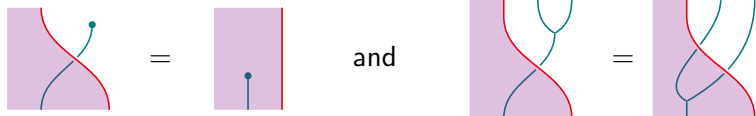
Let  $(\mathbf{X}, T), (\mathbf{Y}, T')$  be monads in  $\mathbf{K}$ .

A *monad opfunctor*  $(F, \phi) : (\mathbf{X}, T) \rightarrow (\mathbf{Y}, T')$  consists of

$F : \mathbf{X} \rightarrow \mathbf{Y}$  and  $\phi : FT \Rightarrow T'F$



such that



# Monads in $\mathbf{Mnd}(\mathbf{K})$

A *monad functor transformation* is a 2-cell  $\sigma : F \Rightarrow F'$  such that



These form a 2-category  $\mathbf{Mnd}^*(\mathbf{K})$ .

When  $\mathbf{X} = \mathbf{Y}$ ,  $T = T'$ , if  $(F, \phi) : (\mathbf{X}, T) \rightarrow (\mathbf{X}, T)$  is a monad, then  $F$  is a monad on  $\mathbf{X}$  and  $\phi$  is a distributive law of  $F$  over  $T$ !

$$\text{i.e.}^2 \mathbf{Dist}(\mathbf{K}) \cong \mathbf{Mnd}^*(\mathbf{Mnd}^*(\mathbf{K}))$$

Also,  $\mathbf{Mnd}^*$  is a monad!

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<sup>2</sup>Can define morphisms between distributive laws such that this is true!

# Algebras over $TS$



# Actions of $T$ , $S$ and $TS$

From before, have monad morphisms<sup>3</sup>:  $T \xrightarrow{T\eta^S} TS \xleftarrow{\eta^TS} S$

$$T\eta^S \begin{array}{c} | \\ \bullet \\ | \end{array} = \begin{array}{c} | \\ | \\ | \end{array} \begin{array}{c} | \\ \bullet \\ | \end{array} ; \quad \eta^TS \begin{array}{c} | \\ \bullet \\ | \end{array} = \begin{array}{c} | \\ \bullet \\ | \end{array} \begin{array}{c} | \\ | \\ | \end{array}$$

These induce  $T$ - and  $S$ -actions on  $U^{TS}$ , via the action of  $TS$ :



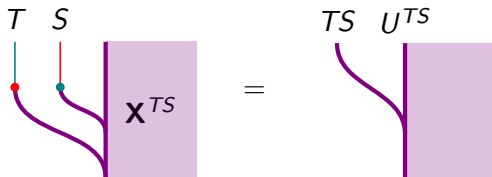
In some sense, any  $TS$ -action is 'captured' by these two actions!

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<sup>3</sup>Monad opfunctors with  $F = 1_x$ .

# Actions of $T$ , $S$ and $TS$

Combining  $T$ - and  $S$ -actions on  $U^{TS}$  gives canonical action of  $TS$ :



Can then show that the  $S$ -action 'distributes over' the  $T$ -action:



# Algebras over $TS$

Let  $\ell$  be a distributive law of  $S$  over  $T$ . From the characterization theorem, we get monads  $TS$  on  $\mathbf{X}$ ,  $\tilde{T}$  on  $\mathbf{X}^S$  and  $\tilde{S}$  on  $\mathbf{X}_T$ .

Theorem (Beck 1969, Cheng 2011)

*The category of algebras of  $TS$  coincides with that of  $\tilde{T}$ .*

$$\mathbf{X}^{TS} \cong (\mathbf{X}^S)^{\tilde{T}}$$

*Dually, the Kleisli category of  $TS$  coincides with that of  $\tilde{S}$ .*

$$\mathbf{X}_{TS} \cong (\mathbf{X}_T)_{\tilde{S}}$$

## Algebras over $TS$

Construct  $\Phi : \mathbf{X}^{TS} \rightarrow (\mathbf{X}^S)^{\tilde{T}}$  and inverse  $\Phi^{-1}$  as lifts arising from universal properties of  $\mathbf{X}^S$ ,  $(\mathbf{X}^S)^{\tilde{T}}$ ,  $\mathbf{X}^{TS}$ :

$$\begin{array}{ccccc} & & (\mathbf{X}^S)^{\tilde{T}} & & \\ & \nearrow \Phi^{-1} & \downarrow U^{\tilde{T}} & \searrow \Phi & \\ \mathbf{X}^{TS} & \dashrightarrow & \mathbf{X}^S & & \mathbf{X}^{TS} \\ & \searrow U^{TS} & \downarrow U^S & \swarrow U^{TS} & \\ & & \mathbf{X} & & \end{array}$$

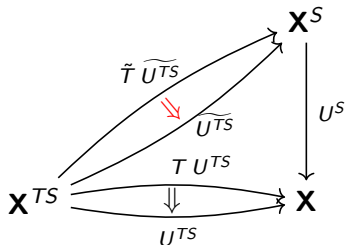
To get  $\Phi^{-1}$ , need  $S$ -action on  $U^{TS}$  and  $\tilde{T}$ -action on lift of  $U^{TS}$ .  
To get  $\Phi$ , need  $TS$  action on  $U^S U^{\tilde{T}}$ .

# Algebras over $TS$

We already have  $T$ - and  $S$ -actions on  $U^{TS}$ .

$S$ -action gives a lift  $\widetilde{U^{TS}} : \mathbf{X}^{TS} \rightarrow \mathbf{X}^S$  of  $U^{TS}$ .

To get  $\widetilde{T}$ -action on  $\widetilde{U^{TS}}$ , lift<sup>4</sup>  $T$ -action on  $U^{TS}$  :



So we have

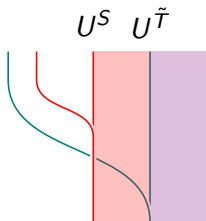
$$\Phi^{-1} : \mathbf{X}^{TS} \rightarrow (\mathbf{X}^S)^{\widetilde{T}}$$

<sup>4</sup>Need  $T$ -action to be an  $S$ -alg. morphism, but this follows from distributivity of  $S$ -action over  $T$ -action.

## Algebras over $TS$

To get  $\Phi : (\mathbf{X}^S)^{\tilde{T}} \rightarrow \mathbf{X}^{TS}$ , need  $TS$ -action on  $U^S U^{\tilde{T}}$ .

Use canonical actions of  $S$  on  $U^S$  and  $\tilde{T}$  on  $U^{\tilde{T}}$ :



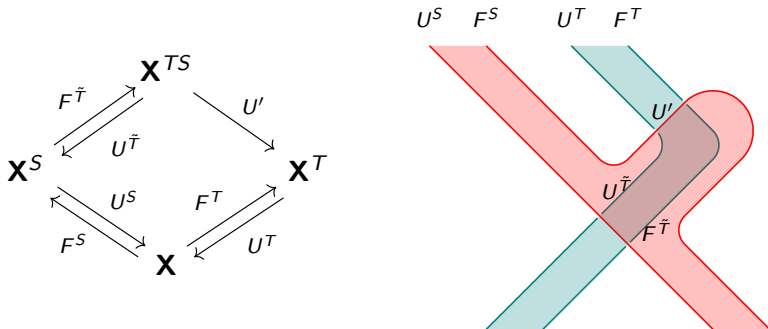
So  $\mathbf{X}^{TS} \cong (\mathbf{X}^S)^{\tilde{T}}$ , and in fact,

$$U^{TS} F^{TS} = TS = U^S U^{\tilde{T}} F^{\tilde{T}} F^S$$

## Distributivity of Adjoints

# Distributivity of adjoints

A distributive law gives rise to a 'distributive square':



where  $U'$  is induced by the  $T$ -action on  $U^{TS}$ .  
 If certain coequalizers exist,  $U'$  has a left adjoint<sup>5</sup>.

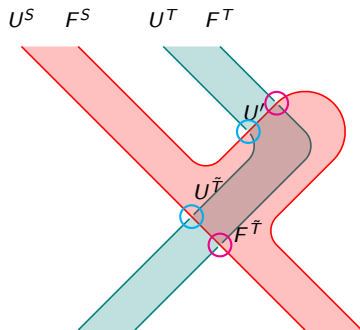
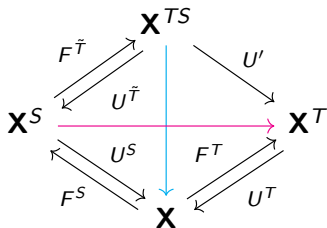
<sup>5</sup>Think of  $U'$  as 'restriction of scalars', and adjoint as 'extension of scalars'.



## Distributivity of adjoints

Both composites  $\mathbf{X}^{TS} \rightarrow \mathbf{X}$  are the same:  $U^S U^{\tilde{T}} = U^T U'$ .

Both composites  $\mathbf{X}^S \rightarrow \mathbf{X}^T$  are the same:  $U' F^{\tilde{T}} = F^T U^S$ .

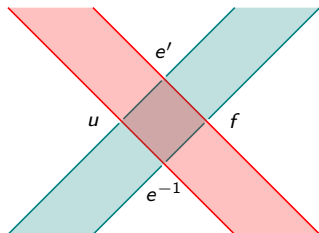
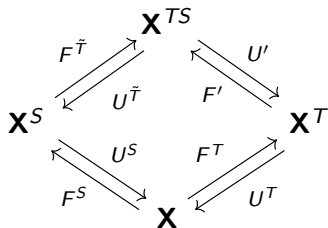


This is a *distributive adjoint situation*, and there is an adjunction:

$$\text{Dist Adj } \mathbf{X} \begin{array}{c} \xrightarrow{\text{Struc}} \\ \perp \\ \xleftarrow{\text{Sem}} \end{array} (\text{Dist } \mathbf{X})^{op}$$

# Distributivity of adjoints

If  $U'$  has an adjoint  $F'$ :



To get distributive law:

Need isomorphisms  $u, f$  that are 'dual' to each other. These give rise to  $e, e'$ .

But  $e$  goes in the 'wrong' direction, so need  $e$  to be an isomorphism too, to get  $e^{-1}$ .

Thank you!

Questions?

## References

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