

# The Enriched Grothendieck Construction

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*In the unenriched world, things are often obscured by all kinds of coincidences.*

- *this talk*

**Main objective:** Present an enriched version of

Theorem (Grothendieck, 1964)

*Let  $B$  be a category. There is a 2-equivalence*

$$\mathbf{Fib}(B) \cong 2\text{-Fun}(B^{op}, \mathbf{Cat})$$

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**Secondary objective:** Highlight things we take for granted in the unenriched setting (i.e. when enriching over **Set**)

- 1 Enriched category theory
- 2 Fibrations
- 3 Results
- 4  $\mathcal{V} = \mathbf{Vect}_k$ ?

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$\otimes$	$\mathbf{1}$	$\mathcal{V}$	$\mathcal{V}$ -categories
$\times$	*	<b>Set</b>	categories
		<b>Cat</b>	strict 2-categories
		<b>sSet</b>	simplicial categories
$\otimes_k$	$k$	<b>Vect<sub>k</sub></b>	$k$ -linear categories
		<b>Ch<sub>R</sub></b>	differential graded categories

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Also, monoids in  $\mathcal{V}$  are  $\mathcal{V}$ -categories with one object:

$$\text{Ob}(\mathcal{C}) = \{*\} \quad \mathcal{C}(*, *) = G, \text{ a monoid}$$

## Underlying categories and free $\mathcal{V}$ -categories

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Conversely, there is often a *free  $\mathcal{V}$ -category*  $\mathcal{C}_{\mathcal{V}}$  on a category  $\mathcal{C}$ .

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One-object example when  $\mathcal{V} = \mathbf{Vect}_k$ :

$\mathcal{C}$	$\mathcal{C}_{\mathcal{V}}$	$(\mathcal{C}_{\mathcal{V}})_0$
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Note:  $\mathcal{C} \neq (\mathcal{C}_{\mathcal{V}})_0$ , but we do have  $\mathcal{C} \hookrightarrow (\mathcal{C}_{\mathcal{V}})_0$ .

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- 2 Fibrations
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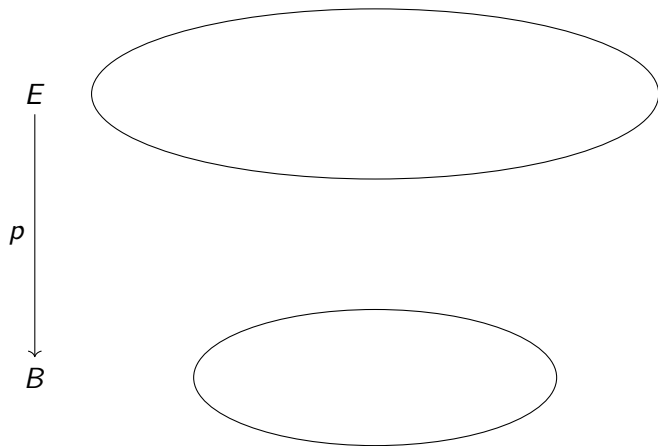
# Fibrations and their fibers

A **fibration** is a functor  $p: E \rightarrow B$  whose fibers are 'nice'.



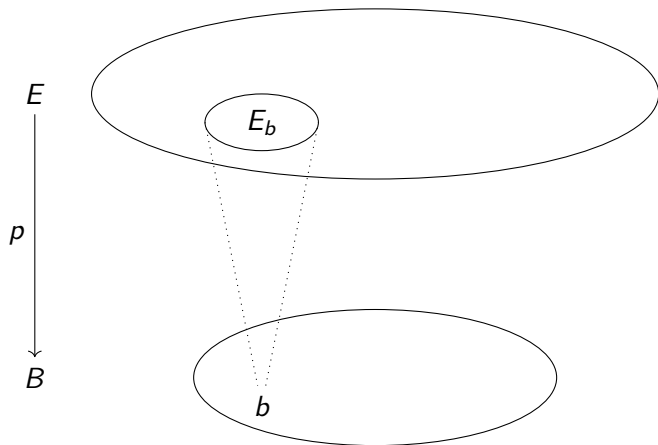
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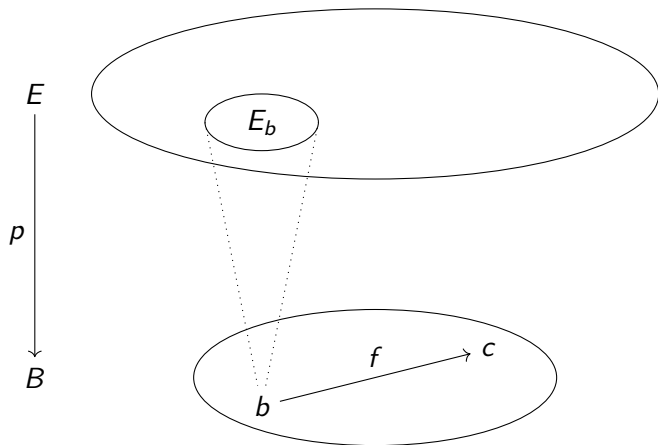
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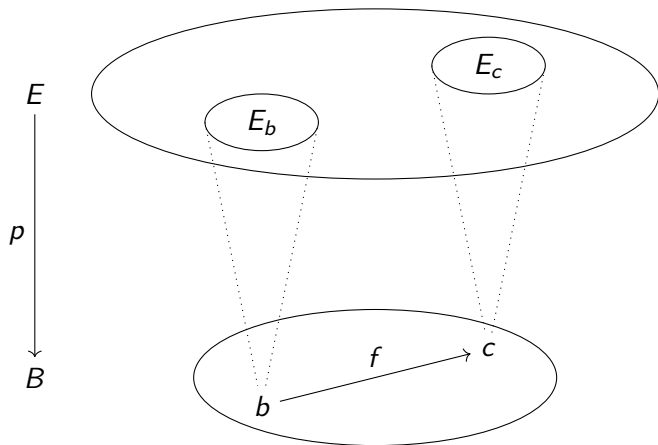
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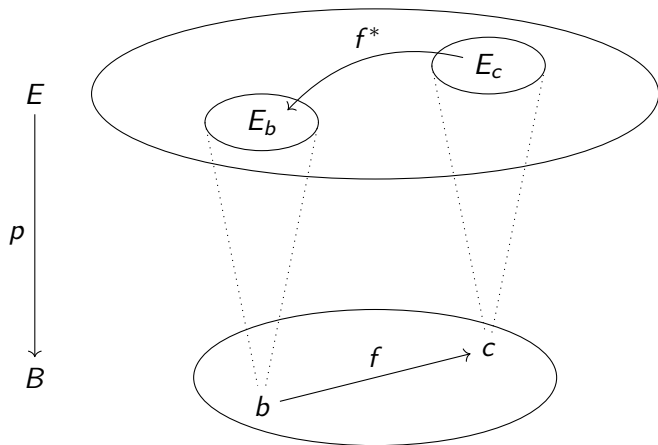
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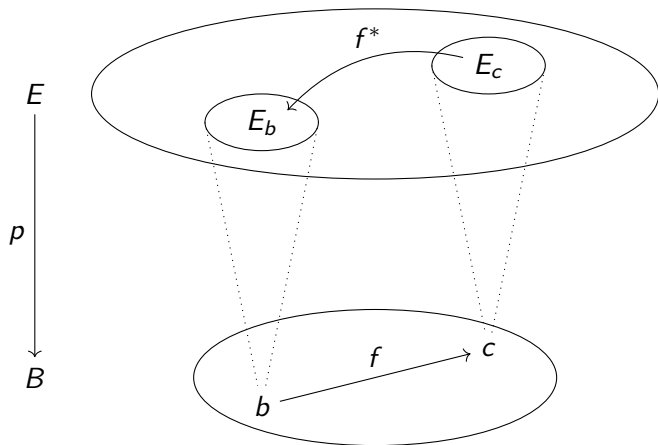
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Further, morphisms in  $E$  are determined by those in the fibers.

# The Grothendieck construction

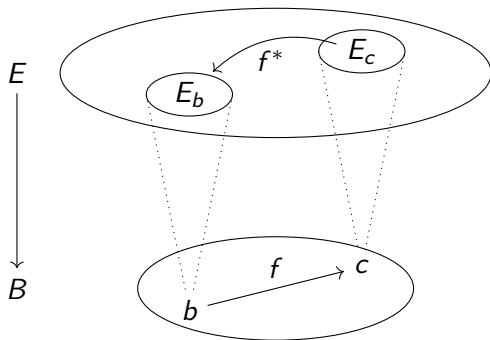
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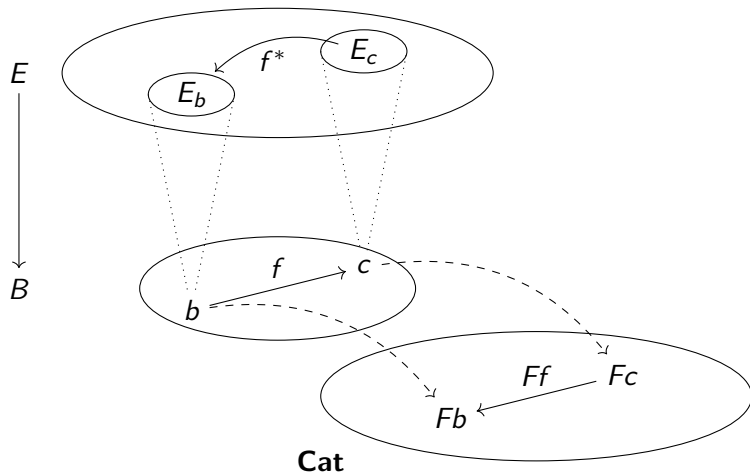




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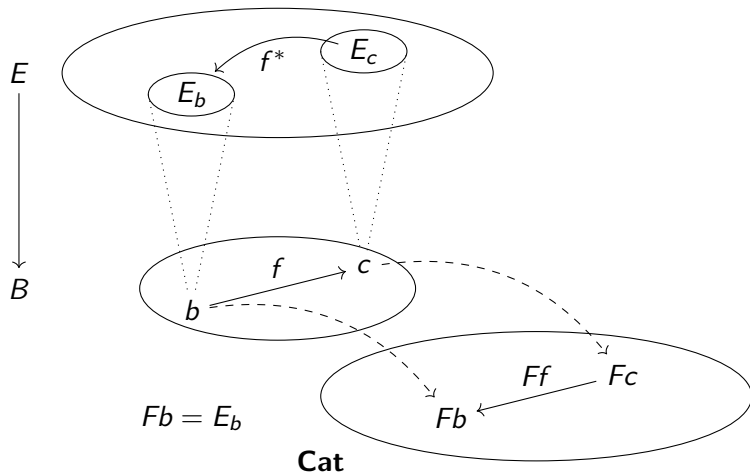
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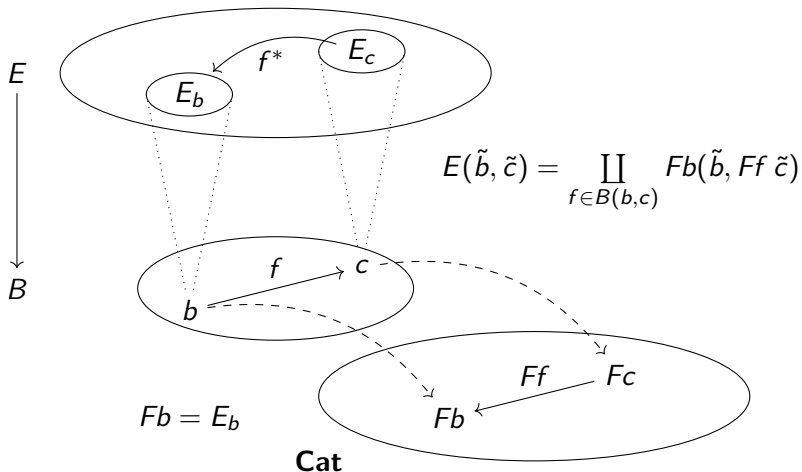
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## Proposition (Beardsley-W.)

Let  $\mathcal{B}$  be a  $\mathcal{V}$ -category. There is a 2-functor

$$\mathcal{V}\text{-Fib}(\mathcal{B}) \rightarrow 2\text{-Fun}(\mathcal{B}_0^{op}, \mathcal{V}\text{-Cat}).$$

## Proposition (Beardsley-W.)

Suppose the unit  $\mathbf{1} \in \mathcal{V}$  is terminal, and pullbacks commute with coproducts in  $\mathcal{V}$ . Let  $B$  be a category. There is 2-functor

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Need to precompose with  $B \hookrightarrow (B_{\mathcal{V}})_0$ .

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When  $\mathcal{V} = \mathbf{sSet}$  [Lurie 2009] or  $\mathbf{Cat}$  [Buckley 2014], there are enhanced equivalences that work for arbitrary  $\mathcal{V}$ -categories  $\mathcal{B}$ .

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- The unit  $k$  is not terminal
- Pullbacks don't commute with coproducts
- Even binary products don't commute with coproducts (both are  $\oplus$ ):

$$X \times (Y \amalg Z) \neq (X \times Y) \amalg (X \times Z)$$

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But  $\otimes_k$  does commute with  $\oplus$ :

$$X \otimes (Y \oplus Z) = (X \otimes Y) \oplus (X \otimes Z)$$

Is there a related category in which  $\otimes$  is a product,  $k$  is terminal, and pullbacks commute with coproducts?

The category  $\mathbf{Coalg}_k = \mathbf{Comon}(\mathbf{Vect}_k)$ :

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But maybe this is the wrong question.

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Maybe substitute some **Sets** by  $\mathcal{V}$  and others by  $\mathbf{Comon}(\mathcal{V})$ :

## Another perspective: comonoids and coactions

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Theorem (Cohen & Montgomery 1984, ..., Tamaki 2009)

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Are fibrations actually special 'coaction functors'?

What other results are secretly about comonoids/coactions?



Thank you!

Questions/comments?